

# Control System of the SPring-8 Storage Ring

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The SPring-8 storage ring is under construction and will be commissioned in February 1997. The design for the control system of the storage ring adopted the so-called Standard Model concept. VMEbus is used as the front-end control system and is widely distributed around the 1436m circumference of the storage ring. The VMEbus is controlled by a single-board-computer with a RISC processor. The VME crate is connected to the FDDI backbone network through an FDDI-Ethernet switching hub. The CPU boards use the LynxOS-based HP-RT real-time operating system. The remote I/O system (RIO), interconnected by optical fibers, is used as the field bus. The operator console system consists of UNIX workstations with a human-oriented interface built using a Motif-based GUI builder. The control software design is based on the event-driven client-server scheme. The remote procedure call (RPC) is used for communication between machine control applications over the network. Accelerator operation history and equipment parameters are stored in a standard database, using an RDBMS.

## 1. Introduction

The SPring-8 project is to construct a large scale advanced synchrotron radiation facility and to promote fundamental science in the field of synchrotron radiation research. The facility being constructed at Harima Science Garden City in western Japan consists of a 1 GeV linac, an 8 GeV booster synchrotron and an 8 GeV low emittance storage ring. The storage ring has 38 beamlines from insertion devices and 23 from bending magnets, including eight beamlines 300m long and three of 1000m. Long beamlines are available for medical experiments planned for the future. Furthermore, SPring-8 is reserving the space of four long straight sections in the ring for future options. The construction of the storage ring is proceeding towards commissioning, which is scheduled for February 1997 [1].

## 2. Storage ring control

### 2.1 Concept

The design concepts of the SPring-8 control system are as follows:

- (1) distributed processors are linked with each other by a high-speed network,
- (2) sub-system (linac, synchrotron, and storage ring) control computers are loosely coupled, due to the different accelerator construction schedules and for the convenience of independent operation,
- (3) all the accelerators are operated from one control room by a small number of operators,
- (4) the system is built using industry standard hardware and software as much as possible,
- (5) for high productivity of application programs, computer-aided program development tools are used.

### 2.2 Architecture

We designed the SPring-8 control system based on the above concept. Figure 1 shows the control system architecture. It consists of UNIX engineering workstations (EWS), VMEbus systems, and optical fiber network (FDDI) systems [2]. In the storage-ring control system the network structure is designed to be flat. The remote I/O systems (RIO) are used mainly as the field bus system [3]. We use HP 9000/700 series EWS with the HP-UX operating system as operator consoles and HP J200 for the database servers. The CPU board of the VME system is HP9000/743rt with a RISC processor. The real-time OS HP-RT was chosen for the CPU board. The control software is designed based on the event-driven client-server scheme. Machine control programs are all written in the C language. The remote procedure call

(ONC/RPC) is used for communication between machine control applications over the network. Accelerator operation history and equipment parameters are stored, using Sybase as the RDBMS. A detailed description is given in the following sections.

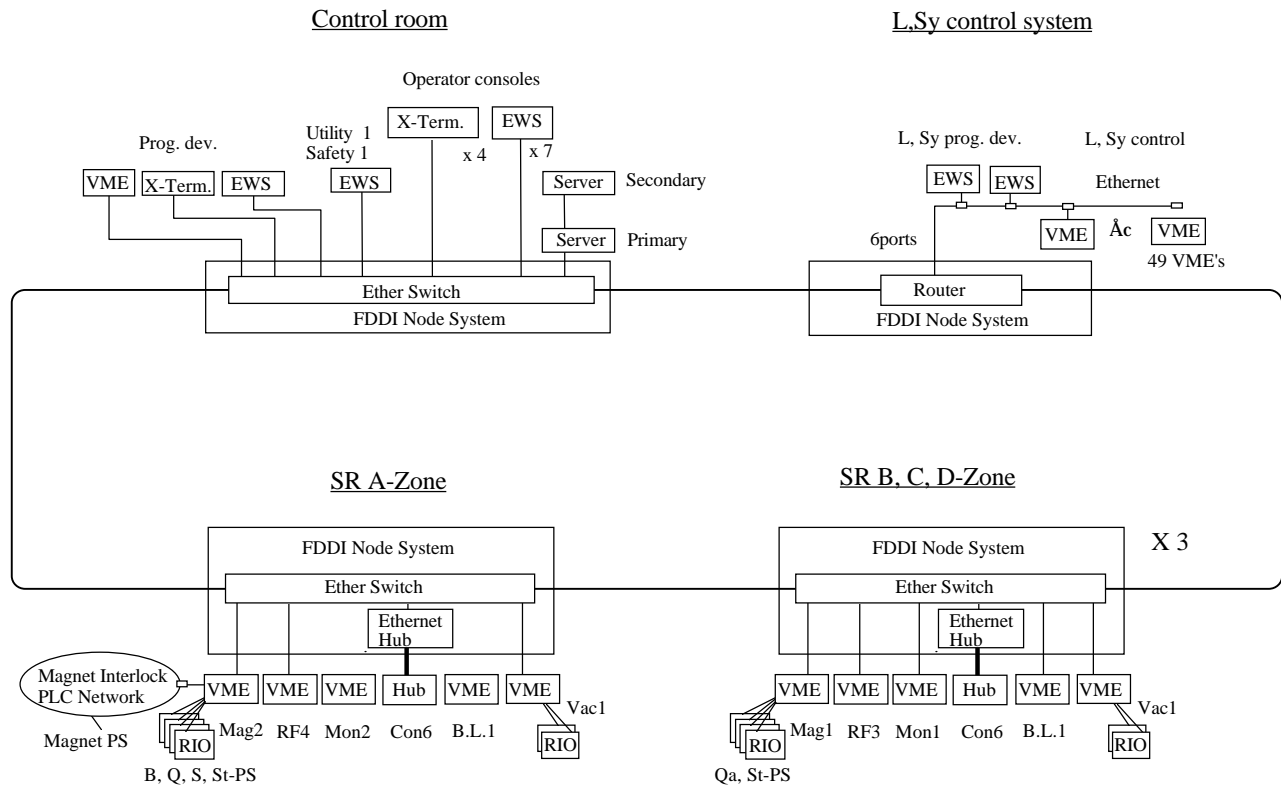


Fig. 1. The architecture of the control system.

### 3. Hardware and software

#### 3.1 VME system

Expandability and maintainability are the most important criteria for adopting industry standard hardware. Consequently the VMEbus system was chosen for the front-end control system, as shown in Fig. 2. The interspersed VME system is suitable for the widely distributed control environment of a large storage ring. The CPU board HP9000/743rt/64, powered by the PA-RISC 7100LC processor, provides quite satisfactory performance (77.7MIPS). It has 16MB of memory and a 20MB PCMCIA flash ROM card, which is used as a boot device for the operating system. VME I/O modules such as DI, DO, TTL DI/O, AI, and PTG are used for direct control of RF equipment. The remote I/O (RIO) system is widely used as the field bus for magnet power supplies, BPMs, and vacuum system controllers. The GPIB bus is used as a field bus as well. We will use 28 VME systems for the storage ring control.

#### 3.2 Front-end operating system

The selection criteria for the OS on the VME system are mainly real-time performance and open system features, indicating compliance with the POSIX standard. From among such OSs HP-RT was chosen. HP-RT is the LynxOS migration to the PA-RISC and can be called a real-time UNIX, so it is easy to understand and manage. The main difference between LynxOS and HP-RT is the program development environment, since HP-RT is supported only in a

cross-development environment. HP-RT is running on the particular CPU boards developed by HP, so although availability maybe limited, we do expect excellent support from the company. This is a crucial point for reducing development cost and time for the new OS. HP-RT device drivers for VME modules are under development by the control group [4].

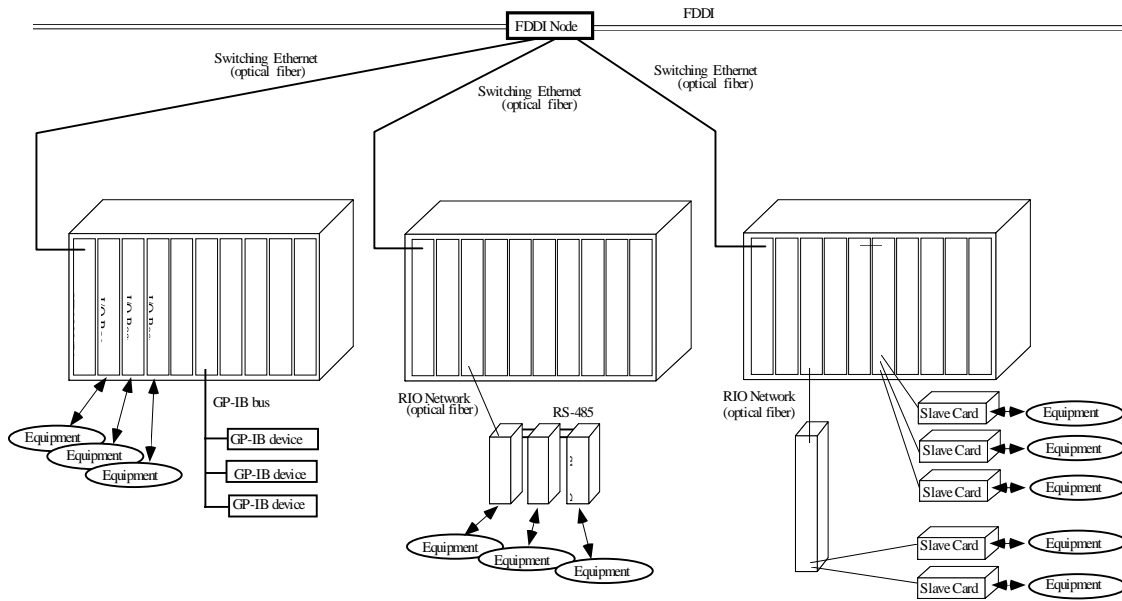


Fig. 2. The VMEbus control system.

### 3.3 Remote I/O bus

An optical fiber distributed remote I/O (RIO) system is used as the field bus for the VME system. The RIO system was initially developed for magnet power-supply control, and eventually will be applied to beam-position monitor and vacuum system control. The RIO system has good electrical isolation and is robust against noise. The system consists of; a) master module with dual port memory, b) several types of slave cards, c) 8-port branch card, and d) star connection with optical fiber cables between master module and slave cards. The features of RIO system are as follows,

- optical fiber connection,
- serial link communication,
- up to 62 slave cards can be connected to one master module,
- 1 Mbps transmission rate,
- up to 1 km transmission distance,
- HDLC protocol,
- can be connected by twisted pair cable (RS-485).

Six types of slave cards are available (type-A, B, D, E, F and G), depending on the type of interface signals of the equipment.

### 3.4 Operator consoles and servers

Operator consoles consist of seven UNIX workstations. These are HP9000/712/60 workstations running HP-UX with 128MB memory and 2GB disk. Four more consoles are planned to be used as the operator consoles for injectors and monitoring of human safety and utilities. Consoles are connected to the FDDI network node through an FDDI-Ethernet switching hub. The GUI based application programs run on each console using the X11 protocol. The connection of consoles has to be considered carefully because unnecessary transactions on the FDDI backbone may deteriorate its performance. In our scheme, switching hubs filter transactions efficiently without leaking unnecessary communication

onto the backbone. The secondary database server is an HP J200 which has dual PA-RISC7200 CPUs, 256MB memory, DAT storage system, CD-ROM juke box and 13GB disk array with RAID mechanism. The primary server will be much faster than the secondary one.

### *3.5 Human-machine interface*

We are using a commercial GUI builder, X-Mate, for developing the operator interface. The look and feel of X-Mate is based on Motif 1.2 with X11 protocol. X-Mate does not necessarily require a window manager such as mwm because it has its own window system on top of Xlib. This is a good point for mission critical 24-hour operation because it is free from

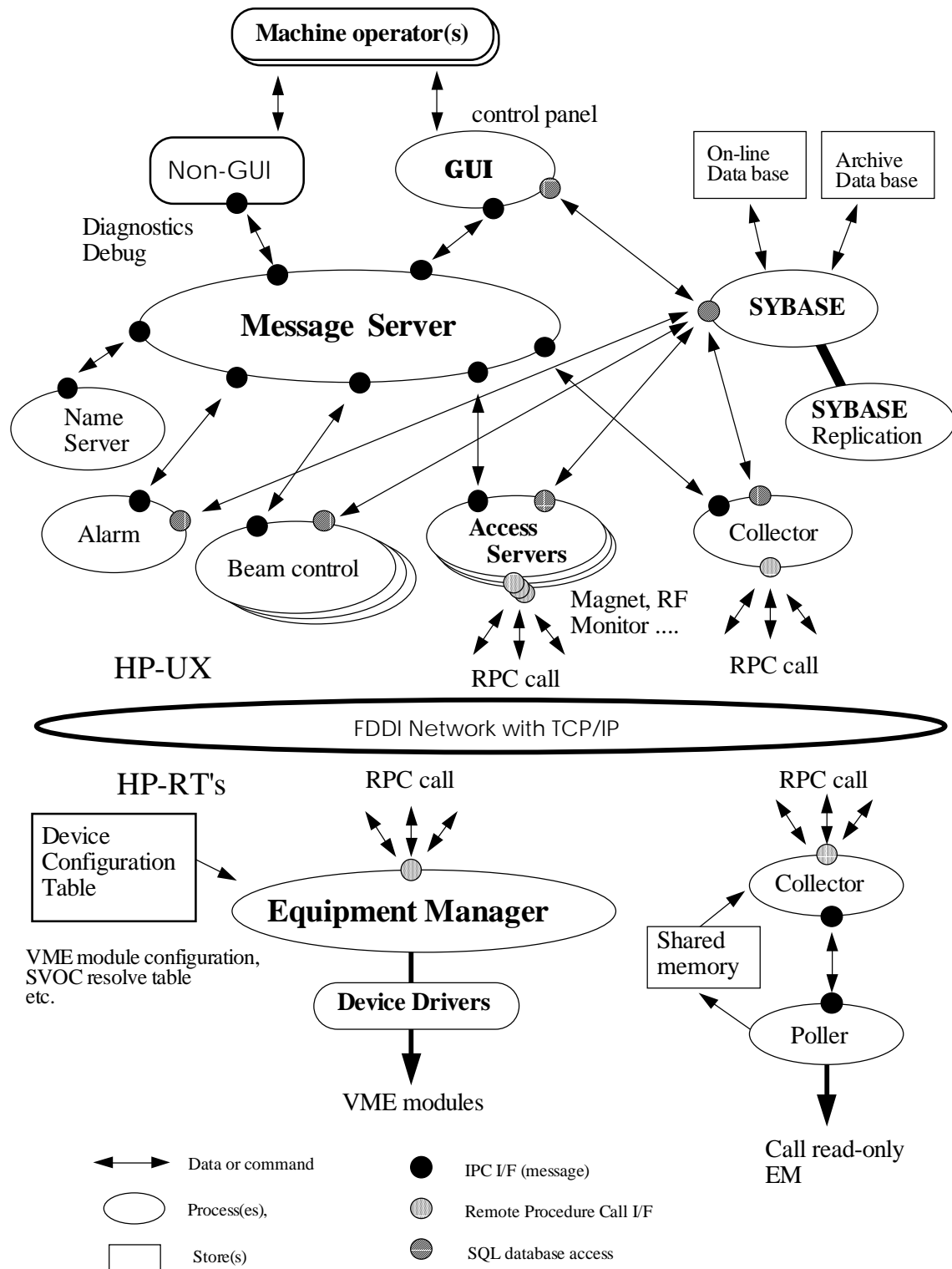


Fig. 3. The software structure of the storage ring control.

the memory leak problem of the Motif window manager. X-Mate provides a good development environment with a comfortable editor. Users are able to make widgets in WYSIWYG without knowing Motif programming. An interface part of the application program has to be written into the call-back routines.

## 4. Equipment access

### 4.1 Access scheme

Figure 3 shows the software scheme for equipment access. An operator command is created using the GUI and is sent out to the Message Server (MS), (see 6.1). The MS will forward the message to the Access Server (AS) after checking its syntax and access control status. The AS parses the message and resolves the destination of the message by asking the Name Server (NS). A composite sequence is analyzed in the AS, and messages are sent to the Equipment Manager (EM) running on the remote CPU board. Each equipment group, e.g. Magnet, RF, Vacuum etc., has its own AS. The synchronization of operations over different equipment is handled here with limited accuracy. The EM is a core process which manages VME devices connected to equipment controllers (see 4.2).

We have an additional cyclic data-acquisition path to get the equipment status. The Poller process running on the CPU board sends pre-registered messages to the read-only EM to get the equipment status. The data is stored in memory by the Poller and taken by the Collector server process on the same board; after collection the data is sent to the Collector client process on the operator console and finally saved to the on-line database by the DBMS, Sybase.

### 4.2 Equipment access

We designed a generic "Equipment Manager" control scheme based on the object-oriented concept [5]. The EM server process running on the front-end CPU board of the VMEbus realizes a device abstraction concept. It controls accelerator components directly without knowing the low-level physical device configuration. The operator commands are resolved by using control semantics and formed into a character string message like an English language syntax S/V/O/C, where the operator action is translated to a verb (V) and controlled equipment is into an object (O) and so on (see 4.3).

Once the EM receives the abstracted command (message), it is translated from logical equipment to physical device(s) by using the translation table. The translation table consists of a device configuration table and the attribute list of equipment. These tables keep the hardware addresses of VME modules and calibration constants etc. Tables are created by an off-line program at the higher level and downloaded to the EM at initialization. After translation, the resolved information of an operator command is passed to the VME device drivers and finally sent to the front-end equipment. The response from the equipment is also abstracted and put into the C of the S/V/O/C format. The reply is sent back to the machine operator through the AS and the MS.

The only way to access the accelerator components is through the EM. The Equipment Manager can be accessed by the Poller locally or by a control application program on the remote consoles by sending a command to the AS in the same manner.

### 4.3 Commands

The S/V/O/C control message consists of the subject (S), verb (V), object (O) and complement (C). The subject has process number, application name, account name, and the host name of the sender which issued the message. The information of the S is used for access control, exclusion control, and security check. The verb is a control action, the object is the equipment to be controlled, and the complement is a value to be written. All accelerator components are abstracted from the physical layer. Programmers of the AS and the Poller are required to know only the abstracted logical object (equipment) itself. The details of control channels and status bits are attached to the object as attributes. The naming scheme for the controlled equipment is chosen to be suitable for beam control at the higher level. Each element and sub-element of the S/V/O/C message is delimited by the characters "/" and "\_". For example, whenever an operator wants to know the current in a storage ring bending magnet, a character string "get/sr\_mag\_ps\_B/current" is the message to be sent. The C returned from the EM can be like "123.45A" or "fail".

## 5. Network

## 5.1 Scheme

As can be seen in Fig. 1, the control system is connected to the duplicated FDDI backbone network with a transmission rate of 100Mbps. There are 6 FDDI nodes for the SPring-8 machine network. One is for the linac and the synchrotron and the others are for the storage ring. All nodes of the storage ring connect to the Ethernet/FDDI switching hub, which provides 10Mbps full bandwidth to each port. We use the optical fiber between the FDDI nodes and the VMEs in order to avoid magnetic interference. The FDDI optical fiber cables are laid around the maintenance corridor of the storage ring and to the injector building. The total length of the cable is about 4500m, and almost all the fibers are laid in a pipe from Air Blown Fiber System (ABFS). The ABFS has multiple pipes inside, and machine control uses one pipe for a bundle of four fibers.

The primary database server will be in the control room and directly connected to FDDI at that node. We will use the replication server system on the secondary database server which will interconnect the machine control network and the beamline user's network. The secondary database server open to users will be a gateway managing access to the control network by machine users.

## 5.2 Protocol

We selected TCP/IP for the network communication protocol and Simple Network Management Protocol (SNMP) for the network management. The TCP/IP protocol is reliable and provides good performance. Reliability is the crucial feature for machine control. The performance of TCP/IP is acceptable for storage ring control and its architecture is suitable because the configuration of our network is relatively simple.

# 6. Process communication

## 6.1 Message distribution

An operator control command is sent by a character string as a message, S/V/O/C as described in 4.1. A message is delivered to the Access Server through the Message Server by using the UNIX System V message mechanism. The MS plays a role as a message distributor. The syntax and access privilege of every message passing over the MS are checked here and all transactions are logged. Each application client running on the workstation has to get a queue number from the MS. Assigned queues are stored onto the Client Registration Table(CRT) managed by the MS. The CRT is used for monitoring the running clients, and its status can be used for surveillance of clients as well. The message distribution scheme is designed to allow only one MS on each console. In the first version, there is no communication between MSs running on different operator consoles; however it may be useful to allow communication between them to enable or disable (exclusion control) clients on different consoles.

## 6.2 Remote communication

The Remote Procedure Call (RPC) is selected for communication over the network because it makes the client-server model more powerful and provides easier programming in the low-level socket interface. There are available commercial implementations of RPCs such as the Sun's Open Network Computing (ONC) and the OSF's Distributed Computing Environment (DCE) based on the HP/Apollo NCS RPC. We selected the ONC/RPC because it is more widely available for the systems which support NFS.

# 7. Database

## 7.1 Management scheme

Data for storage ring control are stored in a unified way and accessed by every process running on consoles in the client-server mode. On-line and off-line applications, distributed in several machines, can store and retrieve the data to/from the database over the network with remote procedure calls (RPC). We chose a commercial relational database management system (RDBMS), Sybase SQL server 10.

Sybase was chosen for its on-line performance as well as its capability as a replication server. A replication server running on server machines can manage the consistency of the data on both the primary and secondary servers. Users of off-line applications, or those which do not affect accelerator operation, are only allowed to access the secondary server. This access control reduces the load on the primary server, which is already heavily loaded by the operational database management. The replication server system plays an important role as a backup server as well.

### *7.2 Accelerator database*

The storage ring database manages three categories of data, constant parameters and on-line and historical data. The parameter database stores relatively stable data of equipment, i.e. calibration constants, physical dimensions, relationships between components. These are managed in the normalized database. The data for status of the accelerator are stored in the on-line database. The Poller/Collector collects current settings and readings of equipment and stores them into the on-line database. The Poller/Collector system is designed to be fast enough to be able to write about 4000 channels of data, which is approximately 10 KB, to the on-line database once per second. The on-line database is denormalized to a flat binary structure to gain performance in updating. The data which are needed for further off-line analysis are partially extracted from the on-line database and stored in the historical database with a time stamp. The archiving cycle of the historical database depends on the data. Alarm status and its history are also written in both the on-line and historical databases.

### *7.3 Data access*

Data in the database are retrieved by issuing an RPC-based SQL request or by using Sybase client-libraries. The higher level processes running on consoles can be SQL clients. On the other hand, processes running on the VME CPU boards can not access the Sybase server directly because that functionality of HP-RT is not available yet. The NFS is one way to access text files on a remote node.

## 8. Status

The design stage of the control hardware has been finished and development is accelerating. The installation of the VME system for magnets and RF has started in the storage ring. A pipe for the air blown fiber system has been installed along the SPring-8 accelerators. We have finished the basic design of the control software, and development of some of the server processes has started using prototype or spiral prototype methods. The conceptual design of the database has started recently.

We have installed the prototype control software on the magnet power supplies and confirmed that the basic control scheme (GUI-MS-AS-EM and Sybase) works successfully on the real system. An operator controlled the power supplies by using the GUI panel, and the data taken by the EM was stored in the Sybase database.

## 9. Plans

The beamline control system is expected to be similar to that of the storage ring in hardware and software, but work has not started yet. The first 10 beamlines, including insertion devices, are planned to be in operation by 1998. It is necessary to finish the beamline control system soon and to fix a clear methodology for control of the insertion devices from the machine control point of view. Digital signal processors (DSP) may be useful for the fast local feedback system for beam orbit correction and image processing of the X-ray beam shape. Recently, the R & D for the DSP system has started.

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